

# CS5201: Advanced Artificial Intelligence

## Prolog programming



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# What is Prolog?

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- Invented early seventies by Alain Colmerauer in France and Robert Kowalski in Britain
- Prolog = Programmation en Logique (Programming in Logic).
- Prolog is a declarative programming language unlike most common programming languages.
- In a declarative language
  - The programmer specifies a goal to be achieved
  - The Prolog system works out how to achieve it
- In purely declarative languages, the programmer only states what the problem is and leaves the rest to the language system

# Relations

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- Prolog programs specify relationships among objects and properties of objects
- When we say, "Ayesha owns the pen", we are declaring the ownership relationship between two objects: Ayesha and the pen.
- When we ask, "Does Ayesha own the pen?" we are trying to find out about a relationship
- Relationships can also be rules such as:
  - Two people are sisters if – they are both female AND they have the same parents.
- In traditional programming relationship may be defined as
  - A and B are sisters if – A and B are both female AND they have the same father AND they have the same mother AND A is not the same as B
- A rule allows us to find out about a relationship even if the relationship is not explicitly stated as a fact

# Programming prolog

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- Declare **facts** describing explicit relationships between objects and properties objects might have (e.g. Subodh likes pizza, sky has\_colour blue)
- Declare **rules** defining implicit relationships between objects and/or rules defining implicit object properties (e.g. X is a parent if there is a Y such that Y is a child of X).
- Then the system can be used by:
  - Asking questions about relationships between objects, and/or about object properties (e.g. does Subodh like pizza? is Ayesha a parent?)

# Prolog & Predicate logic

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- Prolog is a programming language based on predicate logic.
  - A Prolog program attempts to prove a goal, such as `brother(Barney,x)`, from a set of facts and rules.
  - In the process of proving the goal to be true, using substitution and the other rules of inference, Prolog substitutes values for the variables in the goal, thereby "computing" an answer.
- How does Prolog know which facts and which rules to use in the proof?
  - Prolog uses **unification** to determine when two clauses can be made equivalent by a substitution of variables.
  - The unification procedure is used to instantiate the variables in a goal clause based on the facts and rules in the database.

# Tools

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- GNU prolog - `gplc`, `gprolog`
- SWI-Prolog - <https://swi-prolog.org>
- Online tool
  - <https://swish.swi-prolog.org/>
  - [https://www.onlinegdb.com/online\\_prolog\\_compiler](https://www.onlinegdb.com/online_prolog_compiler)

# A simple Prolog program

```
male(albert).
male(edward).
female(alice).
female(victoria).
parent(albert,edward).
parent(victoria,edward).
father(X,Y) :- parent(X,Y), male(X).      %  $\forall X \forall Y ((parent(X, Y) \wedge male(X)) \rightarrow father(X, Y))$ 
mother(X,Y) :- parent(X,Y), female(X).    %  $\forall X \forall Y ((parent(X, Y) \wedge female(X)) \rightarrow mother(X, Y))$ 
```

- A fact/rule (statement) ends with "." and white space ignored
- Read ':'-' after RULE HEAD as "if"
- Read comma in body as "and"
- Comment a line with % or use /\* \*/ for multi-line comments

# Facts & Rules

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- The Prolog environment maintains a set of facts and rules in its database.
  - Facts are axioms; relations between terms that are assumed to be true.
  - Rules are theorems that allow new inferences to be made.

# Facts & Rules

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- The Prolog environment maintains a set of facts and rules in its database.
  - Facts are axioms; relations between terms that are assumed to be true.
  - Rules are theorems that allow new inferences to be made.
- Example facts & rules:
  - `male(adam).`
  - `female(anne).`
  - `parent(adam,barney).`
  - `son(X,Y) :- parent(Y,X) , male(X).`
  - `daughter(X,Y) :- parent(Y,X) , female(X).`
- The first rule is read as follows: for all X and Y, X is the son of Y if there exists X and Y such that Y is the parent of X and X is male.  $\forall X \forall Y ((parent(Y, X) \wedge male(X)) \rightarrow son(X, Y))$
- The second rule is read as follows: for all X and Y, X is the daughter of Y if there exists X and Y such that Y is the parent of X and X is female.  $\forall X \forall Y ((parent(Y, X) \wedge female(X)) \rightarrow daughter(X, Y))$

# Horn clauses

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- To simplify the resolution process in Prolog, statements must be expressed in a simplified form, called Horn clauses.
  - Statements are constructed from terms.
  - Each statement (clause) has (at most) one term on the left hand side of an implication symbol ( :- ).
  - Each statement has a conjunction of zero or more terms on the right hand side.
  - **Example:**  $p_1 \wedge p_2 \wedge \dots \rightarrow q \equiv \neg p_1 \vee \neg p_2 \vee \dots \vee q$
- Prolog has three kinds of statements, corresponding to the structure of the Horn clause used.
  - A **fact** is a clause with an empty right hand side.
  - A **question (or goal)** is a clause with an empty left hand side.
  - A **rule** is a clause with terms on both sides.

# Execution of Prolog program

```
male(albert).
male(edward).
female(alice).
female(victoria).
parent(albert,edward).
parent(victoria,edward).
father(X,Y) :- parent(X,Y), male(X).
mother(X,Y) :- parent(X,Y), female(X).
```

Query:

```
male(X).           %  $\exists X$  male(X)
father(F,edward).  %  $\exists F$  father(F, edward)
father(F,C).       %  $\exists F \exists C$  father(F, C)
```

```
$> gplc family.pl
$> ./family
?- male(albert).
yes
?- male(victoria).
no
?- male(subodh).
no
?- male(X).
X = albert ? ;
X = edward
?- father(F,C).
F=albert, C=edward ? ;
no
```

# Observation about Prolog rules

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- The implication is from right to left
- The scope of a variable is the clause in which it appears.
- Variables whose first appearance is on the left hand side of the clause have implicit universal quantifiers.
- Variables whose first appearance is on the right hand side of the clause have implicit existential quantifiers.
  - $\text{grandparent}(X,Z) \text{ :- } \text{parent}(X,Y), \text{parent}(Y,Z).$
  - $\forall X \forall Z ((\text{parent}(X, Y) \wedge \text{parent}(Y, Z)) \rightarrow \text{grandparent}(X, Z))$

# Basic syntax of Prolog: Terms

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- Constants:
  - Identifiers - sequences of letters, digits, or "\_" that start with lower case letters.
  - Numbers - 1.001
  - Strings enclosed in single quotes
    - Can start with upper case letter or can be a number now treated as a string
- Variables:
  - Sequence of letters digits or "\_" that start with an upper case letter or the \_
    - Underscore by itself is the special "anonymous" variable
- Structures (like function applications)
  - <identifier>(Term-1,...,Term-k)
    - date(20, April, 2020), point(X, Y, Z)
  - Definition can be recursive, so each term can itself be a structure
    - date(+ (5, 15), April, + (2000, - (140, 120)))
  - Structures can be represented as tree

# Arithmetic and logical operators

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- Arithmetic operations: +, -, \*, /
  - The "is" operator forces arithmetic evaluation
  - ?- X is 3 + 8.
- Logical operators: >, <, >=, <=
  - $x =:= y$  (equality check)
  - $x \neq y$  (inequality check)



# Syntax of Prolog: Predicates

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- Predicates are syntactically identical to structured items - <identifier>(Term-1,...,Term-k)
  - male(edward)
  - parent(edward,albert)
  - taller\_than(subodh,shyam)
  - likes(X)
    - Note that X is a variable. X can take on any term as value so that this fact asserts
- Facts make assertion

# Syntax of Prolog: Facts and Rules

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- Rules: `PredicateH :- predicate-1, ..., predicate-k.`
  - First predicate is *rule head*. Terminated by a period
  - Rules encode ways of deriving or computing a new fact
  - `animal(X) :- elephant(X).` - X can be concluded to be animal if it shown that X is elephant
  - `taller_than(X,Y) :- height(X,H1), height(Y,H2), H1 > H2.`
  - `father(X,Y) :- parent(X,Y), male(X).`

# Operation of Prolog

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- A query is a sequence of predicates: predicate-1, predicate-2, ..., predicate-k
- Prolog tries to prove that this sequence of predicates is true using the facts and rules in the Prolog Program.
- In proving the sequence it performs the computation you want.
- Example:
  - elephant(fred).
  - elephant(mary).
  - elephant(joe).
  - animal(fred) :- elephant(fred).
  - animal(mary) :- elephant(mary).
  - animal(joe) :- elephant(joe).
  - QUERY: animal(fred), animal(mary), animal(joe).

# Operation

---

- Starting with the first predicate P1 of the query Prolog examines the program from TOP to BOTTOM.
- It finds the first RULE HEAD or FACT that matches P1
- Then it replaces P1 with the RULE BODY.
- If P1 is matched a FACT, we can think of FACTs as having empty bodies (so P1 is simply removed).
- The result is a new query.
- Example
  - P1 :- Q1, Q2, Q3.
  - QUERY: P1, P2, P3.
  - P1 matches with the rule, therefore, new QUERY: Q1, Q2, Q3, P2, P3.

# Execution of Prolog program

```
elephant(fred).  
elephant(mary).  
elephant(joe).  
animal(fred) :- elephant(fred).  
animal(mary) :- elephant(mary).  
animal(joe) :- elephant(joe).  
QUERY: animal(fred), animal(mary), animal(joe).
```

1. elephant(fred), animal(mary), animal(joe).
2. animal(mary), animal(joe).
3. elephant(mary), animal(joe).
4. animal(joe).
5. elephant(joe).
6. EMPTY QUERY

# Operation

---

- If this process reduces the query to the empty query, Prolog returns "yes".
- However, during this process each predicate in the query might match more than one fact or rule head
  - Prolog always choose the first match it finds. Then if the resulting query reduction did not succeed (i.e., we hit a predicate in the query that does not match any rule head of fact), Prolog backtracks and tries a new match.

# Execution of Prolog program

---

```
ant_eater(fred).  
animal(fred) :- elephant(fred).  
animal(fred) :- ant_eater(fred).  
QUERY: animal(fred)
```

1. elephant(fred).
2. FAIL BACKTRACK
3. ant\_eater(fred).
4. EMPTY QUERY

# Operation

---

- Backtracking can occur at every stage as the query is processed

p(1) :- a(1).

p(1) :- b(1).

a(1) :- c(1).

c(1) :- d(1).

c(1) :- d(2).

b(1) :- e(1).

e(1).

d(3).

QUERY: p(1)

p(1)

# Operation

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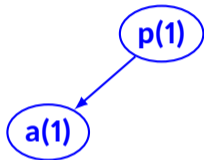
c(1) :- d(2).

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e(1).

d(3).

QUERY: p(1)



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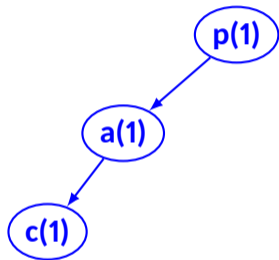
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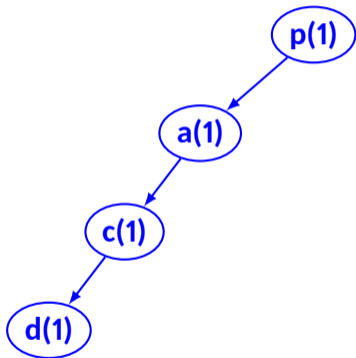
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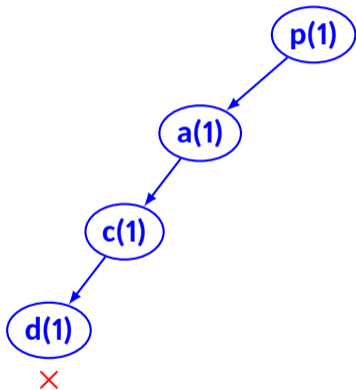
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d(3).

QUERY: p(1)



# Operation

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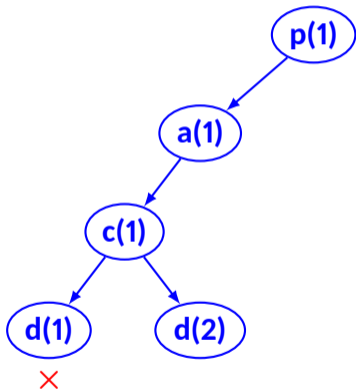
$c(1) :- d(2).$

$b(1) :- e(1).$

$e(1).$

$d(3).$

QUERY:  $p(1)$



# Operation

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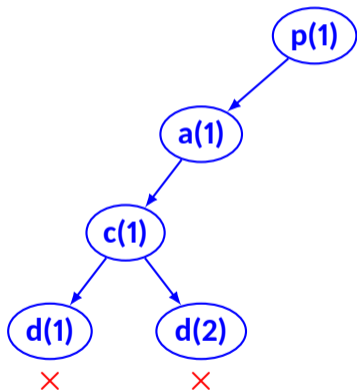
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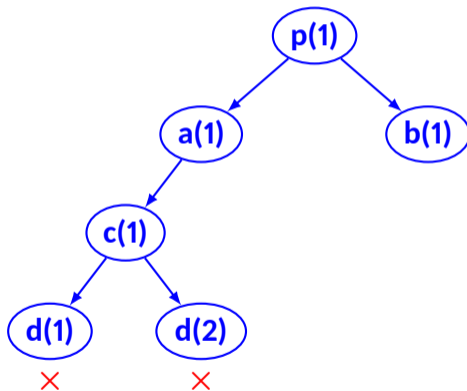
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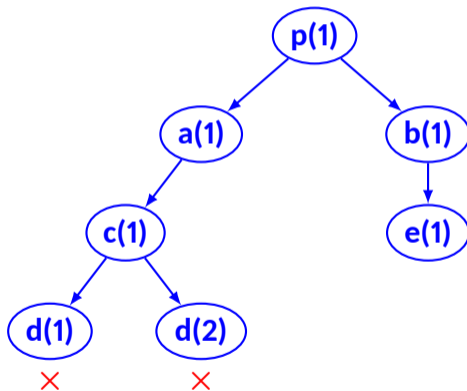
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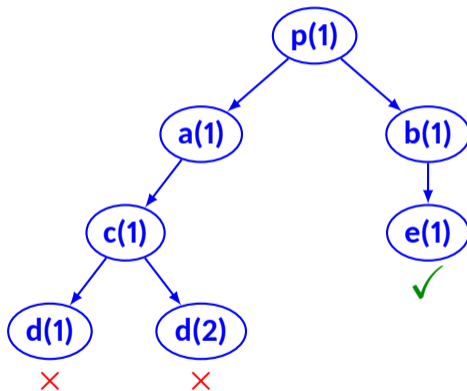
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e(1).

d(3).

QUERY: p(1)



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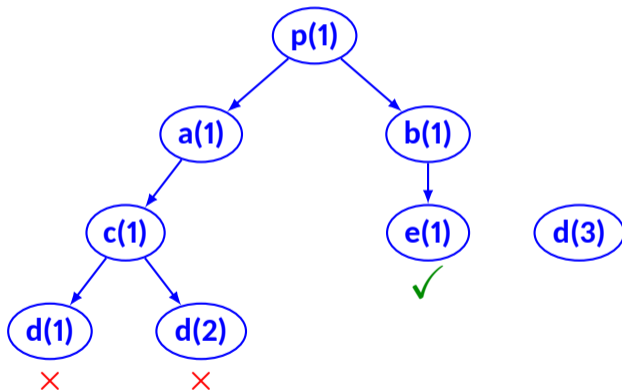
c(1) :- d(2).

b(1) :- e(1).

e(1).

d(3).

QUERY: p(1)



# Operation

- With backtracking we can get more answers by using ";"

p(1) :- a(1).

p(1) :- b(1).

a(1) :- c(1).

c(1) :- d(1).

c(1) :- d(2).

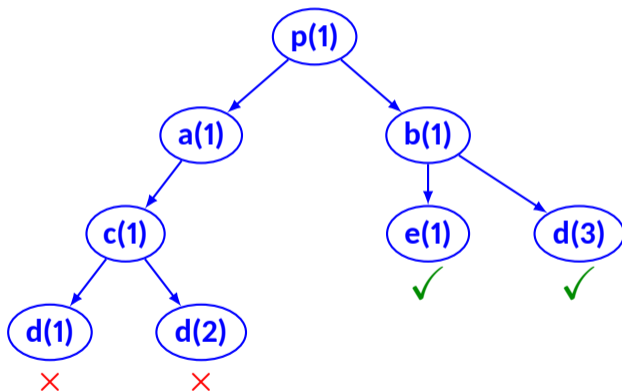
b(1) :- e(1).

**b(1) :- d(3).**

e(1).

d(3).

QUERY: p(1)



# Variables and Matching

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- Variables allow us to
  - Compute more than yes/no answer, compress the program
- Example:
  - elephant(fred).
  - elephant(mary).
  - elephant(joe).
  - animal(fred) :- elephant(fred).
  - animal(mary) :- elephant(mary).
  - animal(joe) :- elephant(joe).
- The three rules can be replaced by the single rule animal(X) :- elephant(X).
- When matching queries against rule heads (of facts) variables allow many additional matches.

# Example

---

```
elephant(fred).  
elephant(mary).  
elephant(joe).  
animal(X) :- elephant(X).  
QUERY: animal(fred), animal(mary), animal(joe)
```

1. X=fred, elephant(X), animal(mary), animal(joe)
2. animal(mary), animal(joe)
3. X=mary, elephant(X), animal(joe)
4. animal(joe)
5. X=joe, elephant(X)
6. EMPTY QUERY

# Operation with Variables

---

- Queries are processed as before (via rule and fact matching and backtracking), but now we can use variables to help us match rule heads or facts.
- A query predicate matches a rule head or fact (either one with variables) if
  - The predicate name must match. So `elephant(X)` can match `elephant(joe)`, but can never match `ant_eater(joe)`.
  - Then, the arity of the predicates are matched (number of terms). So `foo(X,Y)` can match `foo(joe,mary)`, but cannot match `foo(joe)` or `foo(joe,mary,fred)`.
  - If the predicate names and arities match then each of the k-terms match. So for `foo(t1, t2, t3)` to match `foo(s1, s2, s3)` we must have that `t1` matches `s1`, `t2` matches `s2`, and `t3` matches `s3`.
  - During this matching process we might have to "bind" some of the variables to make the terms match.
  - These bindings are then passed on into the new query (consisting of the rule body and the left over query predicates).

# Variable matching (Unification)

---

- Two terms with variables match if :
  - If both are constants (identifiers, numbers, or strings) and are identical
  - If one or both are bound variables then they match if what the variables are bound to match
    - X and mary where X is bound to the value mary will match
    - X and Y where X is bound to mary and Y is bound to mary will match
    - X and ann where X is bound to mary will not match
  - If one of the terms is an unbound variable then they match and we bind the variable to the term
    - X and mary where X is unbound match and make X bound to mary.
    - X and Y where X is unbound and Y is bound to mary match and make X bound to mary.
    - X and Y where both X and Y are unbound match and make X bound to Y (or vice versa).

# Solving queries

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- Prolog work as follows
  - Unification
  - Goal directed reasoning
  - Rule ordering
  - DFS and backtracking

# List processing in Prolog

---

- Much of prolog's computation is organized around lists. Two key things we do with a list is iterate over them and build new ones.
- Checking membership: `member(X,Y)` - X is a member of list Y

# List processing in Prolog

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  - `member(X,[X|_])`.
  - `member(X,[_|T]) :- member(X,T)`.

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  - `member(X,[_|T]) :- member(X,T)`.
- Building a list of integers in range [i,j] (`build(from, to, NewList)`)
  - `build(I,J,[]) :- I > J`.
  - `build(I,J,[I | Rest]) :- I =< J, N is I+1, build(N,J,Rest)`.

# List examples

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- Concatenation: `concatenation(X, Y, Z)`

# List examples

---

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  - `concatenation([], L, L)`.
  - `concatenation([X|L1], L2, [X|L3]) :- concatenation(L1, L2, L3)`.

# List examples

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- Example:
  - `concatenation([a,b], [c,d], Y)`.

# List examples

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  - `concatenation([X|L1], L2, [X|L3]) :- concatenation(L1, L2, L3)`.
- Example:
  - `concatenation([a,b], [c,d], Y)`.
  - `X=a, concatenation([X|b], [c,d], [X|Y1])`.
  - `concatenation([b], [c,d], Y1)`.
  - `X=b, concatenation([X|[]], [c,d], [X|Y2])`.
  - `concatenation([], [c,d], Y2)`.

# List examples

---

- Adding in front: `add(X, Y, Z)` – add X in front of Y and result into Z
  - `add(X, L, [X|L])`.

# List examples

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- Adding in front: `add(X, Y, Z)` – add X in front of Y and result into Z
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- Deletion: `del(X, Y, Z)` – delete X from Y and store result in Z

# List examples

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  - `del(X, [X|Tail], Tail)`.
  - `del(X, [Y|Tail], [Y|Tail1]) :- del(X, Tail, Tail1)`.

# List examples

---

- Adding in front: `add(X, Y, Z)` – add X in front of Y and result into Z
  - `add(X, L, [X|L]).`
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  - `del(X, [X|Tail], Tail).`
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- Sublist: `sublist(S, L)` – check whether S is sublist of L

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  - `sublist(S, L) :- concatenation(L1, L2, L), concatenation(S, L3, L2)`.

# List examples

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- Adding in front: `add(X, Y, Z)` – add X in front of Y and result into Z
  - `add(X, L, [X|L]).`
- Deletion: `del(X, Y, Z)` – delete X from Y and store result in Z
  - `del(X, [X|Tail], Tail).`
  - `del(X, [Y|Tail], [Y|Tail1]) :- del(X, Tail, Tail1).`
- Sublist: `sublist(S, L)` – check whether S is sublist of L
  - `sublist(S, L) :- concatenation(L1, L2, L), concatenation(S, L3, L2).`
- GCD: `gcd(X, Y, Z)`
  - `gcd(X,X,X).`
  - `gcd(X,Y,Z) :- X < Y, Y1 is Y-X, gcd(X,Y1,Z).`

# List examples

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- **Permutation:**
  - `permutation([], []).`
  - `permutation([X|L], P) :- permutation(L, L1), insert(X, L1, P).`

# 8-queens

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- **Solution:**

- `solution(Queens):- permutation([1,2,3,4,5,6,7,8], Queens), safe(Queens).`

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  - `permutation([Head | Tail], PermList) :-  
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  - `safe([]).`
  - `safe([Queen | Others]) :- safe(Others), noattack(Queen, Others, 1).`

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  - `solution(Queens):- permutation([1,2,3,4,5,6,7,8], Queens), safe(Queens).`
- **Permutation:**
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  - `permutation([Head | Tail], PermList) :-  
permutation(Tail, Permtail), del(Head, PermList, Permtail).`
- **Safe:**
  - `safe([]).`
  - `safe([Queen | Others]) :- safe(Others), noattack(Queen, Others, 1).`
- **No-attack:**
  - `noattack(_,[],_).`
  - `noattack(Y,[Y1 | Ylist],Xdist) :-  
Y-Y1 =\= Xdist, Y1-Y =\= Xdist, Dist is Xdist+1, noattack(Y, Ylist, Dist).`

# Source of inefficiency

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- $f(X, \text{normal}) :- X < 3.$
- $f(X, \text{alert1}) :- 3 \leq X, X < 6.$
- $f(X, \text{alert2}) :- 6 \leq X.$
- Query1:  $f(2, Y), Y = \text{alert1}.$

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- Query2:  $f(7, Y).$

# Source of inefficiency

---

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- $f(X, \text{normal}) :- X < 3, !.$
- $f(X, \text{alert1}) :- 3 \leq X, X < 6, !.$
- $f(X, \text{alert2}) :- 6 \leq X.$
- Query2:  $f(7, Y).$

## Efficient version

- $f(X, \text{normal}) :- X < 3, !.$
- $f(X, \text{alert1}) :- X < 6, !.$
- $f(X, \text{alert2}) :- 6 \leq X.$

# Cuts – controlled backtracking

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$C :- P, Q, R, !, S, T, U.$

$C :- V.$

$A :- B, C, D.$

?-  $A.$

- Backtracking within the goal list  $P, Q, R$
- As soon as the cut is reached:
  - All alternatives of  $P, Q, R$  are suppressed
  - The clause  $C :- V$  will also be discarded
  - Backtracking is possible within  $S, T, U$ .
  - No effect for rule  $A$ , that is backtracking within  $B, C, D$  remains active.

# Cuts

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- `del_duplicates([], [])`.
- `del_duplicates([Head | Tail], Result) :-`  
    `member(Head, Tail), !, del_duplicates(Tail, Result). % R1`
- `del_duplicates([Head | Tail], [Head | Result]) :- del_duplicates(Tail, Result). % R2`

# Cuts - Example

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- If  $X \geq Y$  then  $\text{Max}=X$ , otherwise  $\text{Max}=Y$ 
  - `max(X, Y, X) :- X >= Y, !.`
  - `max(X, Y, Y).`
- Adding an element into a list without duplication
  - `add(X, L, L) :- member(X, L), !.`
  - `add(X, L, [X|L]).`

*Thank you!*