CS5201: Advanced Artificial Intelligence

State Space Search



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Complex problems & solutions-1







$$\int \frac{\mathbf{x}^4}{(1-\mathbf{x}^2)^{\frac{5}{2}}} d\mathbf{x}$$

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Automated problem solving by search

- Generalized techniques for solving large class of complex problems
- Problem statement is the input and solution is the output (sometime problem specific algorithm / method could be the output)
- Al search based problem formulation requires following steps broadly
 - Configuration or state
 - Constraints or definitions of valid configuration
 - Rules for change of state and their outcomes
 - Initial state or start configurations
 - Goal satisfying configurations
 - An implicit state space
 - Valid solutions from start to goal in that state space
 - General algorithms which search for solutions in this state space
- Challenges: Size of implicit state space, Capturing domain knowledge, Intelligent algorithms that work in reasonable time and memory, Handling incompleteness and uncertainty

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- Start and goal states

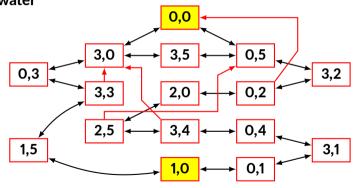
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 - Depending on the problem formulation, it can be a PATH from Start to Goal or a Sub-graph of And-ed Nodes
- Search algorithms
 - Intelligently explore the Implicit Graph or State Space by examining only a small sub-set to find the solution
 - To use Domain Knowledge or HEURISTICS to try and reach Goals faster

Two jug problem

• There is a large bucket B full of water and Two (O2) jugs, J1 of volume 3 liter and J2 of volume 5 liter. You are allowed to fill up any empty jug from the bucket, pour all water back to the bucket from a jug or pour from one jug to another. The goal is to have jug J1 with exactly one (O1) liter of water

- State-space modeling
 - State definition: (J1,J2)
 - Rules:
 - Fill(J1), Fill(J2)
 - Empty(J1), Empty(J2)
 - Pour(J1,J2), Pour(J2,J1)
 - Start state: (0,0)
 - Goal state: (1,0)



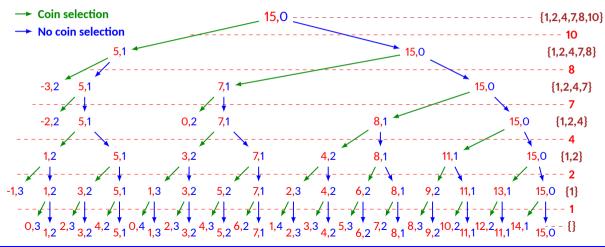
Partial state-space

Coin change: State-space

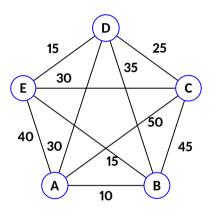
- Given a set of coins C, what is the minimum number coins required to provide sum S?
- Example: $C = \{1, 2, 4, 7, 8, 10\}, S = 15$

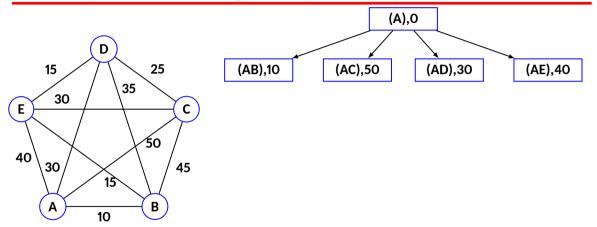
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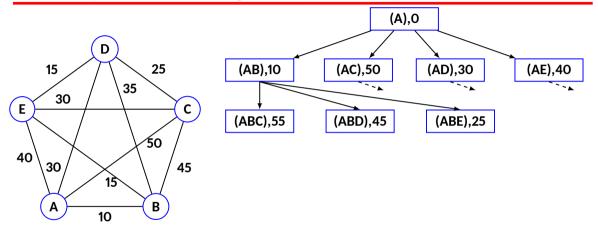
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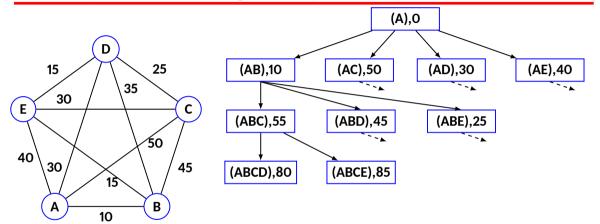


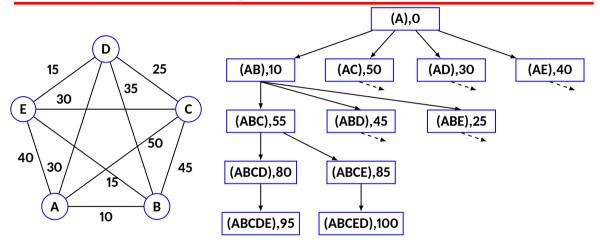
(A),0

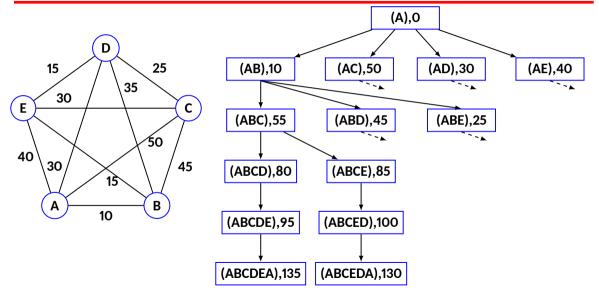








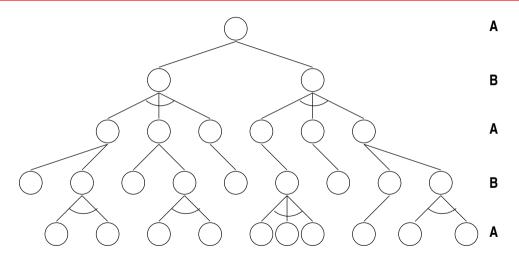




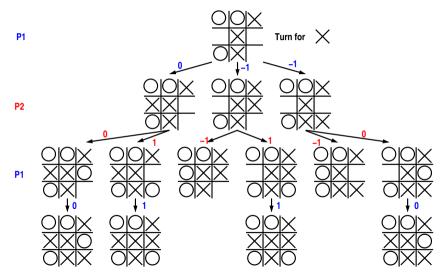
Modelling AND/OR graphs

- OR nodes are ones for which one has a choice
- The AND nodes could be compositional (sum, product, min, max, etc., depending on the subproblems are composed)
 - Adversarial where the other parties have choice, usually in games
 - Probabilistic environmental actions

AND/OR graphs



Adversarial AND/OR graphs



Compositional AND/OR graph

- Let A and B be two matrices of size $p \times q$ and $q \times r$. Therefore, to determine $A \times B$ we need to perform $p \times q \times r$ number of multiplications
- Let A_1 be a 10×20 matrix, A_2 be a 20×5 and A_3 be a 5×50 matrix
- Then the number of computation
 - $(A_1 \times A_2) \times A_3$

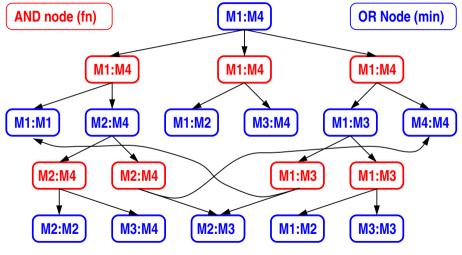
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 - $(A_1 \times A_2) \times A_3 (10 \times 20 \times 5) + (10 \times 5 \times 50) = 3500$
 - $A_1 \times (A_2 \times A_3)$

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 - $(A_1 \times A_2) \times A_3 (10 \times 20 \times 5) + (10 \times 5 \times 50) = 3500$
 - $A_1 \times (A_2 \times A_3) (20 \times 5 \times 50) + (10 \times 20 \times 50) = 15000$

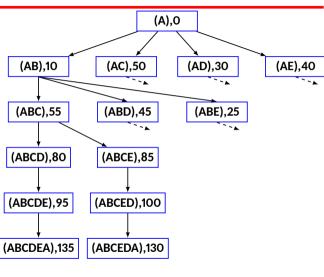
MCM: AND/OR graphs



M1(M2(M3M4)) = (M1M2)(M3M4) = ((M1M2)M3))M4 = (M1(M2M3))M4 = M1((M2M3)M4)

Searching implicit graph

- Given the start state the SEARCH Algorithm will create successors based on the State Transformation Rules and make part of the Graph EXPLICIT.
- It will EXPAND the Explicit graph IN-TELLIGENTLY to rapidly search for a solution without exploring the entire Implicit Graph or State Space
- For OR Graphs, the solution is a PATH from start to Goal.
- Cost is usually sum of the edge costs on the path, though it could be something based on the problem



Searching implicit graph: Algorithms

- Basic algorithms Depth-First (DFS), Breadth-First (BFS), Iterative deepening (IDS)
- Cost-based algorithms Depth-First Branch-and-Bound, Best First Search, Best-First Iterative deepening
- Widely used algorithms A* and IDA* (OR graphs), AO* (AND/OR graphs), Alphabeta pruning (Game-trees)

Basic algorithms in OR graphs: DFS

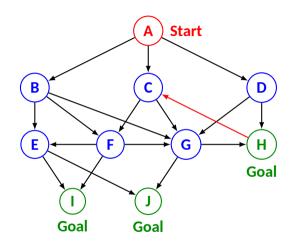
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- 2. [Select] Select the first Node *n* on the OPEN List. If OPEN is empty, Terminate
- 3. [Goal Test] If *n* is Goal, then decide on Termination or Continuation / Cost Updation
- 4. [Expand]
 - a. Generate the successors n_1, n_2, \ldots, n_k , of node n_k based on the State Transformation Rules
 - **b.** Put n in CLOSED List
 - c. For each n_i , not already in OPEN or CLOSED List, put n_i in the FRONT of OPEN List
 - d. For each n_i already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2

Basic algorithms in OR graphs: IDS

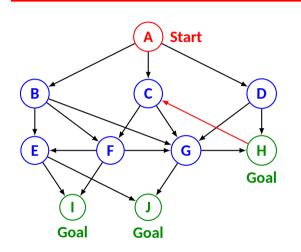
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 - **b.** Put *n* in LIST CLOSED
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 - d. For each n_i already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2
- IDS performs DFS level wise manner (DFS(1), DFS(2),)

Basic algorithms in OR graphs: BFS

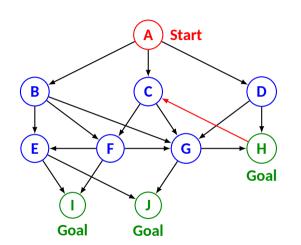
- 1. [Initialize] Initially the OPEN List contains the Start Node s. CLOSED List is Empty.
- 2. [Select] Select the first Node *n* on the OPEN List. If OPEN is empty, Terminate
- 3. [Goal Test] If *n* is Goal, then decide on Termination or Continuation / Cost Updation
- 4. [Expand]
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 - **b.** Put *n* in LIST CLOSED
 - c. For each n_i , not already in OPEN or CLOSED List, put n_i in the END of OPEN List
 - d. For each n_i already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2



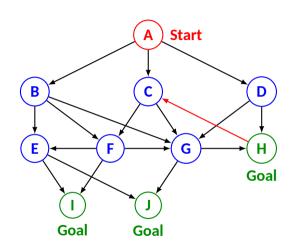
- Depth-first search
- Breadth-first search
- Iterative deepening search



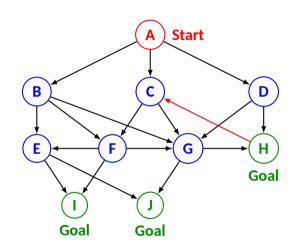
Step OPEN CLOSED



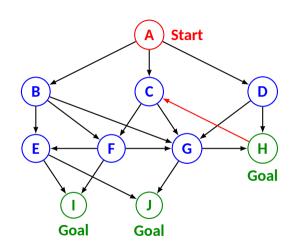
Step	OPEN	CLOSED
1	Α	{}



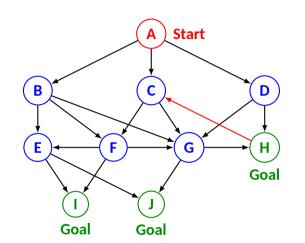
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α



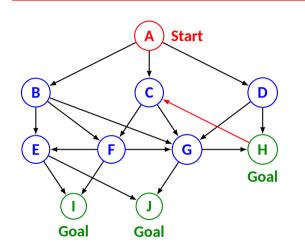
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	EFGCD	AB



Step	OPEN	CLOSED
1	Α	{}
2	BCD	A
3	EFGCD	AB
4	IJFGCD	ABE



Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	EFGCD	AB
4	IJFGCD	ABE
5	I is goal no	de, can terminate

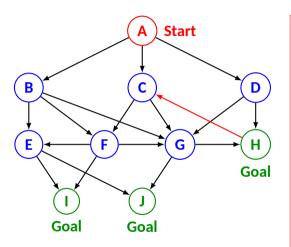


Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	EFGCD	AB
4	IJFGCD	ABE
5	I is goal node, can terminate	

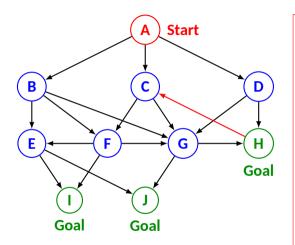
Search can continue for more goal nodes if minimum length or cost is a criteria.

DFS may not terminate if there is an infinite depth path even if there is a goal node at finite depth.

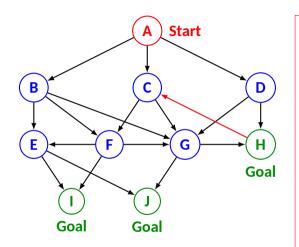
DFS has low memory overhead.



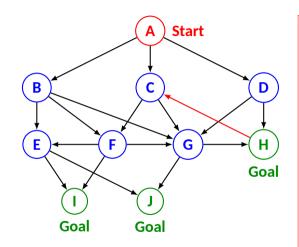
Outcome Step



Step	Outcome
1	DFS(L=1) - No solution



Step	Outcome
1	DFS(L=1) - No solution
2	DFS(L=2) - Goal node H reached



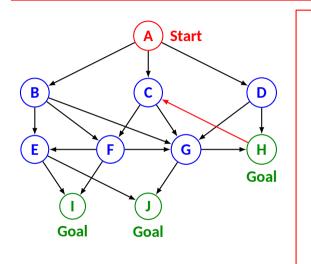
Step	Outcome
1	DFS(L=1) - No solution
2	DFS(L=2) - Goal node H reached
3	Can terminate with path from A to
	Н

This is guaranteed to be the minimum length path.

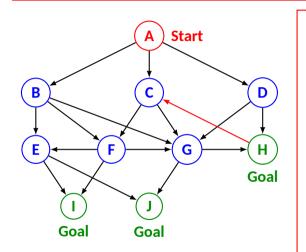
IDS guarantees shortest length path to goal.

IDS may re-expand nodes many times.

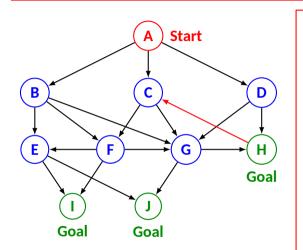
IDS has lower memory requirement than BFS.



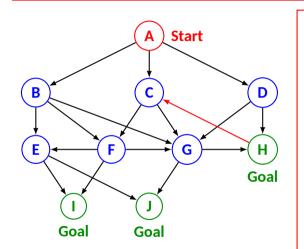
Step **OPEN CLOSED**



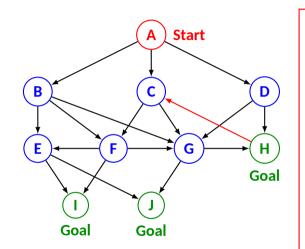
Step	OPEN	CLOSED
1	Α	{}



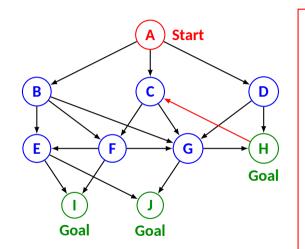
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α



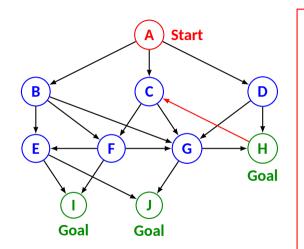
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB



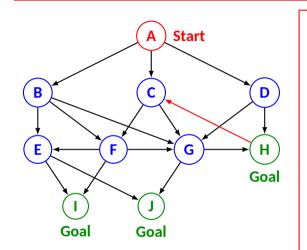
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC



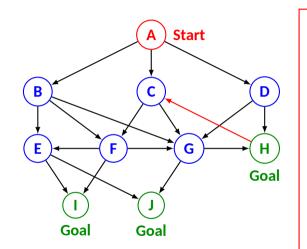
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD



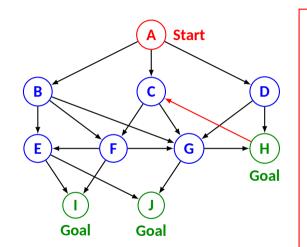
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE



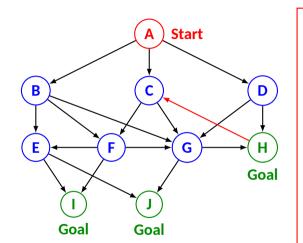
Step	OPEN	CLOSED
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4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE
7	GHIJ	ABCDEF



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3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE
7	GHIJ	ABCDEF
8	HIJ	ABCDEFG



Step	OPEN	CLOSED
1	Α	{}
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3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE
7	GHIJ	ABCDEF
8	HIJ	ABCDEFG
9	Goal no	de H found



Step	OPEN	CLOSED	
1	Α	{}	
2	BCD	Α	
3	CDEFG	AB	
4	DEFG	ABC	
5	EFGH	ABCD	
6	FGHIJ	ABCDE	
7	GHIJ	ABCDEF	
8	HIJ	ABCDEFG	
9	Goal node H found		

BFS guarantees shortest length path to goal but has higher memory requirement.

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	Iterative Deepening	BFS
Complete?			

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Criterion	DFS	Iterative Deepening	BFS
Complete?	No	Yes	Yes

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Criterion	DFS	Iterative Deepening	BFS
Complete?	No	Yes	Yes
Time			

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- Optimality under the assumption of identical cost for all steps

Criterion	DFS	Iterative Deepening	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$

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Criterion	DFS	Iterative Deepening	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space			

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Criterion	DFS	Iterative Deepening	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space	O(bm)	O(bd)	$O(b^d)$

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- Optimality under the assumption of identical cost for all steps

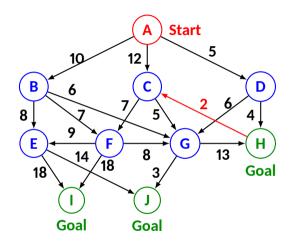
Criterion	DFS	Iterative Deepening	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space	O(bm)	O(bd)	$O(b^d)$
0			

Optimal

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	Iterative Deepening	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space	O(bm)	O(bd)	$O(b^d)$
Optimal	No	Yes	Yes

Searching state space graph with edge cost

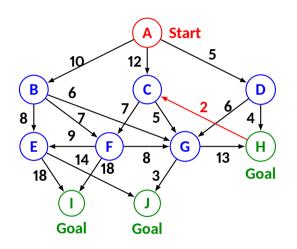


Modifying basic algorithms to incorporate cost

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 - d. For each n_i already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2

Algorithm IDS Performs DFS level by level iteratively (DFS(1), DFS(2), ...and so on)

Searching state space graph with edge cost



Cost ordered search:

- DFBB
- Best first search
- Best first IDS
- Use of heuristic estimates: A*, AO*

Thank you!