# **CS5201: Advanced Artificial Intelligence**

# **State Space Search**



#### **Arijit Mondal**

Dept of Computer Science and Engineering Indian Institute of Technology Patna www.iitp.ac.in/~arijit/

## **Complex problems & solutions-1**







$$\int \frac{\mathbf{x}^4}{(1-\mathbf{x}^2)^{\frac{5}{2}}} d\mathbf{x}$$

CS5201, SPRING, 2025



## **Automated problem solving by search**

- Generalized techniques for solving large class of complex problems
- Problem statement is the input and solution is the output (sometime problem specific algorithm / method could be the output)
- Al search based problem formulation requires following steps broadly
  - Configuration or state
  - Constraints or definitions of valid configuration
  - Rules for change of state and their outcomes
  - Initial state or start configurations
  - Goal satisfying configurations
  - An implicit state space
  - Valid solutions from start to goal in that state space
  - General algorithms which search for solutions in this state space
- Challenges: Size of implicit state space, Capturing domain knowledge, Intelligent algorithms that work in reasonable time and memory, Handling incompleteness and uncertainty

- State or configurations
  - A set of variables which define a state or configuration
  - Domains for every variable and constraints among variables to define a valid configuration

- State or configurations
  - A set of variables which define a state or configuration
  - Domains for every variable and constraints among variables to define a valid configuration
- State transformation rules or moves
  - A set of rules which define which are the valid set of next state of a given state
  - It also indicates who can make these moves (OR Nodes, AND nodes, etc)

- State or configurations
  - A set of variables which define a state or configuration
  - Domains for every variable and constraints among variables to define a valid configuration
- State transformation rules or moves
  - A set of rules which define which are the valid set of next state of a given state
  - It also indicates who can make these moves (OR Nodes, AND nodes, etc)
- State space or implicit graph
  - The Complete Graph produced out of the state transformation rules.
  - Typically too large to store. Could be Infinite.

- State or configurations
  - A set of variables which define a state or configuration
  - Domains for every variable and constraints among variables to define a valid configuration
- State transformation rules or moves
  - A set of rules which define which are the valid set of next state of a given state
  - It also indicates who can make these moves (OR Nodes, AND nodes, etc)
- State space or implicit graph
  - The Complete Graph produced out of the state transformation rules.
  - Typically too large to store. Could be Infinite.
- Start and goal states

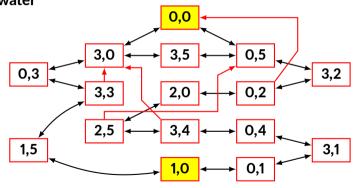
- State or configurations
  - A set of variables which define a state or configuration
  - Domains for every variable and constraints among variables to define a valid configuration
- State transformation rules or moves
  - A set of rules which define which are the valid set of next state of a given state
  - It also indicates who can make these moves (OR Nodes, AND nodes, etc)
- State space or implicit graph
  - The Complete Graph produced out of the state transformation rules.
  - Typically too large to store. Could be Infinite.
- Start and goal states
- Solutions, costs
  - Depending on the problem formulation, it can be a PATH from Start to Goal or a Sub-graph of And-ed Nodes

- State or configurations
  - A set of variables which define a state or configuration
  - Domains for every variable and constraints among variables to define a valid configuration
- State transformation rules or moves
  - A set of rules which define which are the valid set of next state of a given state
  - It also indicates who can make these moves (OR Nodes, AND nodes, etc)
- State space or implicit graph
  - The Complete Graph produced out of the state transformation rules.
  - Typically too large to store. Could be Infinite.
- Start and goal states
- Solutions, costs
  - Depending on the problem formulation, it can be a PATH from Start to Goal or a Sub-graph of And-ed Nodes
- Search algorithms
  - Intelligently explore the Implicit Graph or State Space by examining only a small sub-set to find the solution
  - To use Domain Knowledge or HEURISTICS to try and reach Goals faster

## Two jug problem

• There is a large bucket B full of water and Two (O2) jugs, J1 of volume 3 liter and J2 of volume 5 liter. You are allowed to fill up any empty jug from the bucket, pour all water back to the bucket from a jug or pour from one jug to another. The goal is to have jug J1 with exactly one (O1) liter of water

- State-space modeling
  - State definition: (J1,J2)
  - Rules:
    - Fill(J1), Fill(J2)
    - Empty(J1), Empty(J2)
    - Pour(J1,J2), Pour(J2,J1)
  - Start state: (0,0)
  - Goal state: (1,0)



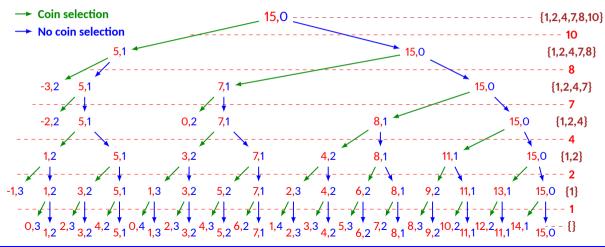
Partial state-space

## Coin change: State-space

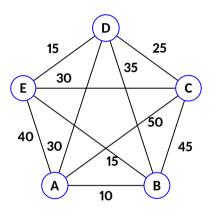
- Given a set of coins C, what is the minimum number coins required to provide sum S?
- Example:  $C = \{1, 2, 4, 7, 8, 10\}, S = 15$

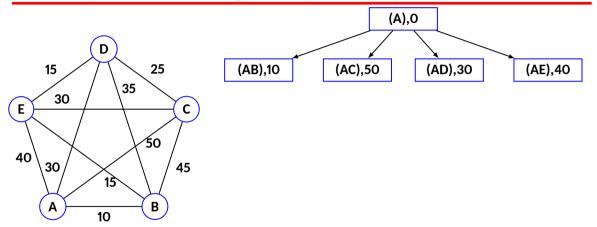
#### **Coin change: State-space**

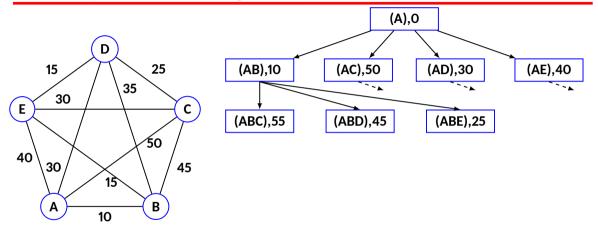
- Given a set of coins C, what is the minimum number coins required to provide sum S?
- Example:  $C = \{1, 2, 4, 7, 8, 10\}, S = 15$

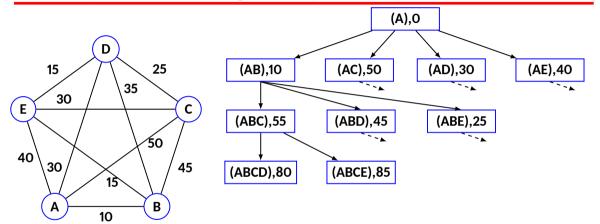


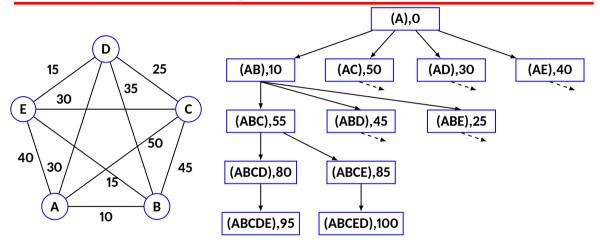
(A),0

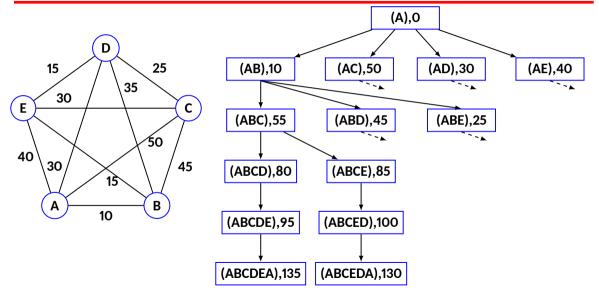








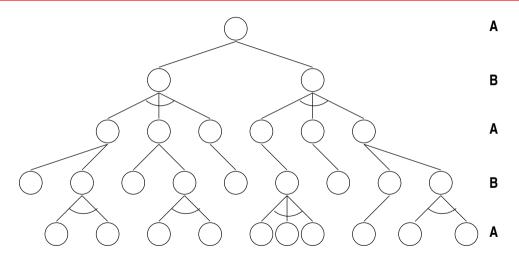




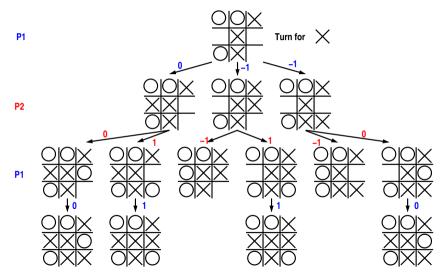
### Modelling AND/OR graphs

- OR nodes are ones for which one has a choice
- The AND nodes could be compositional (sum, product, min, max, etc., depending on the subproblems are composed)
  - Adversarial where the other parties have choice, usually in games
  - Probabilistic environmental actions

## **AND/OR graphs**



## **Adversarial AND/OR graphs**



## **Compositional AND/OR graph**

- Let A and B be two matrices of size  $p \times q$  and  $q \times r$ . Therefore, to determine  $A \times B$  we need to perform  $p \times q \times r$  number of multiplications
- Let  $A_1$  be a  $10 \times 20$  matrix,  $A_2$  be a  $20 \times 5$  and  $A_3$  be a  $5 \times 50$  matrix
- Then the number of computation
  - $(A_1 \times A_2) \times A_3$

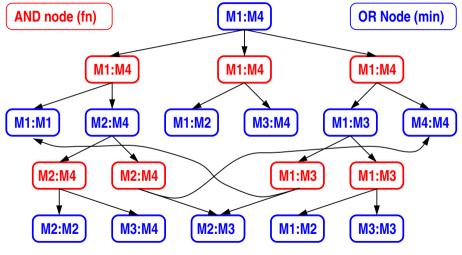
## **Compositional AND/OR graph**

- Let A and B be two matrices of size  $p \times q$  and  $q \times r$ . Therefore, to determine  $A \times B$  we need to perform  $p \times q \times r$  number of multiplications
- Let  $A_1$  be a  $10 \times 20$  matrix,  $A_2$  be a  $20 \times 5$  and  $A_3$  be a  $5 \times 50$  matrix
- Then the number of computation
  - $(A_1 \times A_2) \times A_3 (10 \times 20 \times 5) + (10 \times 5 \times 50) = 3500$
  - $A_1 \times (A_2 \times A_3)$

## **Compositional AND/OR graph**

- Let A and B be two matrices of size  $p \times q$  and  $q \times r$ . Therefore, to determine  $A \times B$  we need to perform  $p \times q \times r$  number of multiplications
- Let  $A_1$  be a  $10 \times 20$  matrix,  $A_2$  be a  $20 \times 5$  and  $A_3$  be a  $5 \times 50$  matrix
- Then the number of computation
  - $(A_1 \times A_2) \times A_3 (10 \times 20 \times 5) + (10 \times 5 \times 50) = 3500$
  - $A_1 \times (A_2 \times A_3) (20 \times 5 \times 50) + (10 \times 20 \times 50) = 15000$

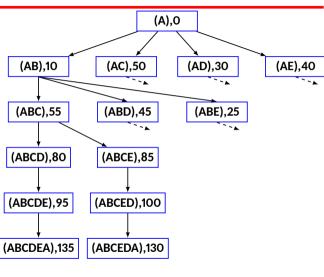
#### MCM: AND/OR graphs



M1(M2(M3M4)) = (M1M2)(M3M4) = ((M1M2)M3))M4 = (M1(M2M3))M4 = M1((M2M3)M4)

### Searching implicit graph

- Given the start state the SEARCH Algorithm will create successors based on the State Transformation Rules and make part of the Graph EXPLICIT.
- It will EXPAND the Explicit graph IN-TELLIGENTLY to rapidly search for a solution without exploring the entire Implicit Graph or State Space
- For OR Graphs, the solution is a PATH from start to Goal.
- Cost is usually sum of the edge costs on the path, though it could be something based on the problem



## **Searching implicit graph: Algorithms**

- Basic algorithms Depth-First (DFS), Breadth-First (BFS), Iterative deepening (IDS)
- Cost-based algorithms Depth-First Branch-and-Bound, Best First Search, Best-First Iterative deepening
- Widely used algorithms A\* and IDA\* (OR graphs), AO\* (AND/OR graphs), Alphabeta pruning (Game-trees)

## **Basic algorithms in OR graphs: DFS**

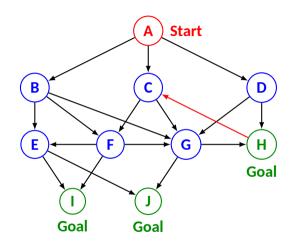
- 1. [Initialize] Initially the OPEN List contains the Start Node s. CLOSED List is Empty.
- 2. [Select] Select the first Node *n* on the OPEN List. If OPEN is empty, Terminate
- 3. [Goal Test] If *n* is Goal, then decide on Termination or Continuation / Cost Updation
- 4. [Expand]
  - a. Generate the successors  $n_1, n_2, \ldots, n_k$ , of node  $n_k$  based on the State Transformation Rules
  - **b.** Put n in CLOSED List
  - c. For each  $n_i$ , not already in OPEN or CLOSED List, put  $n_i$  in the FRONT of OPEN List
  - d. For each  $n_i$  already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2

## **Basic algorithms in OR graphs: IDS**

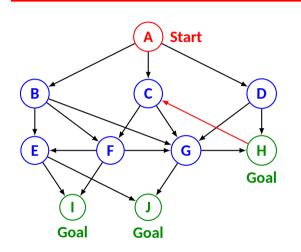
- 1. [Initialize] Initially the OPEN List contains the Start Node s. CLOSED List is Empty.
- 2. [Select] Select the first Node *n* on the OPEN List. If OPEN is empty, Terminate
- 3. [Goal Test] If *n* is Goal, then decide on Termination or Continuation / Cost Updation
- 4. [Expand]
  - a. Generate the successors  $n_1, n_2, \ldots, n_k$ , of node n, based on the State Transformation Rules
  - **b.** Put *n* in LIST CLOSED
  - c. For each  $n_i$ , not already in OPEN or CLOSED List, put  $n_i$  in the FRONT of OPEN List
  - d. For each  $n_i$  already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2
- IDS performs DFS level wise manner (DFS(1), DFS(2), ....)

## **Basic algorithms in OR graphs: BFS**

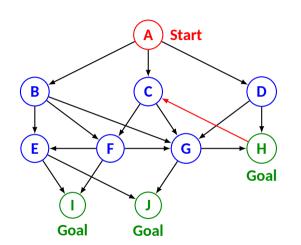
- 1. [Initialize] Initially the OPEN List contains the Start Node s. CLOSED List is Empty.
- 2. [Select] Select the first Node *n* on the OPEN List. If OPEN is empty, Terminate
- 3. [Goal Test] If *n* is Goal, then decide on Termination or Continuation / Cost Updation
- 4. [Expand]
  - a. Generate the successors  $n_1, n_2, \ldots, n_k$ , of node  $n_k$  based on the State Transformation Rules
  - **b.** Put *n* in LIST CLOSED
  - c. For each  $n_i$ , not already in OPEN or CLOSED List, put  $n_i$  in the END of OPEN List
  - d. For each  $n_i$  already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2



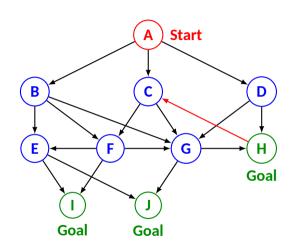
- Depth-first search
- Breadth-first search
- Iterative deepening search



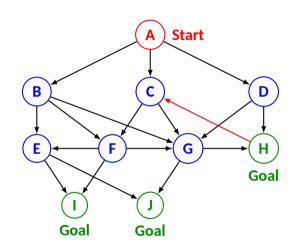
**Step OPEN CLOSED** 



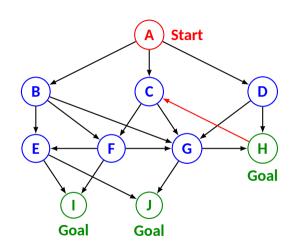
Step	OPEN	CLOSED
1	Α	{}



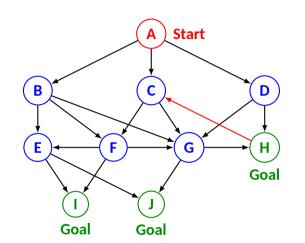
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α



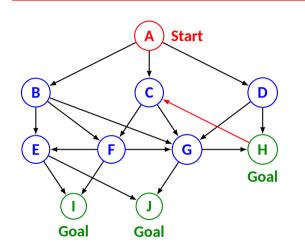
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	EFGCD	AB



Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	EFGCD	AB
4	IJFGCD	ABE



Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	EFGCD	AB
4	IJFGCD	ABE
5	I is goal no	de, can terminate

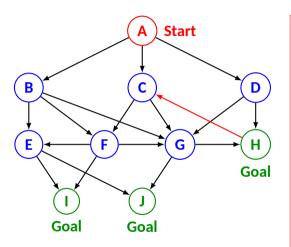


Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	EFGCD	AB
4	IJFGCD	ABE
5	I is goal node, can terminate	

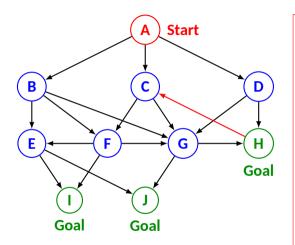
Search can continue for more goal nodes if minimum length or cost is a criteria.

DFS may not terminate if there is an infinite depth path even if there is a goal node at finite depth.

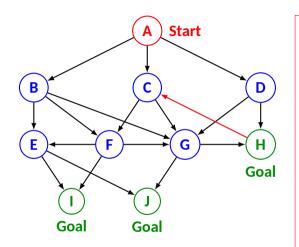
DFS has low memory overhead.



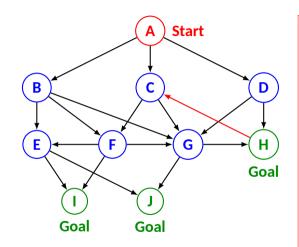
**Outcome** Step



Step	Outcome
1	DFS(L=1) - No solution



Step	Outcome
1	DFS(L=1) - No solution
2	DFS(L=2) - Goal node H reached



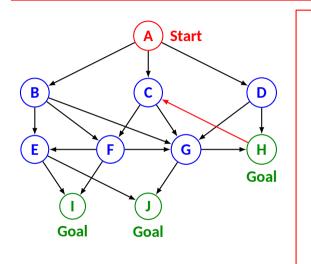
Step	Outcome
1	DFS(L=1) - No solution
2	DFS(L=2) - Goal node H reached
3	Can terminate with path from A to
	Н

This is guaranteed to be the minimum length path.

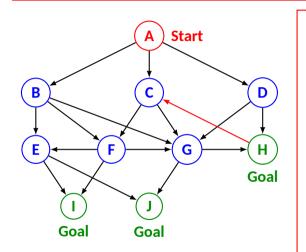
IDS guarantees shortest length path to goal.

IDS may re-expand nodes many times.

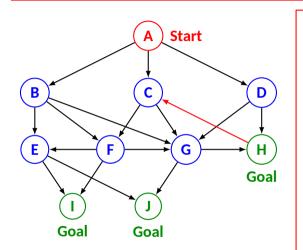
IDS has lower memory requirement than BFS.



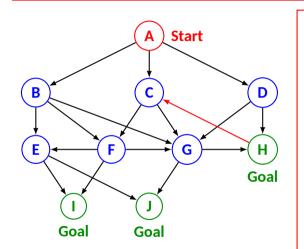
Step **OPEN CLOSED** 



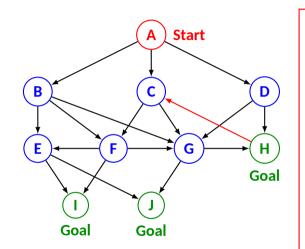
Step	OPEN	CLOSED
1	Α	{}



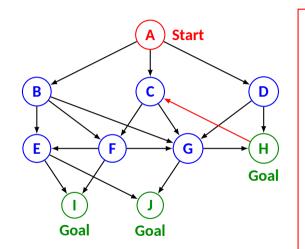
Step	<b>OPEN</b>	CLOSED
1	Α	{}
2	BCD	Α



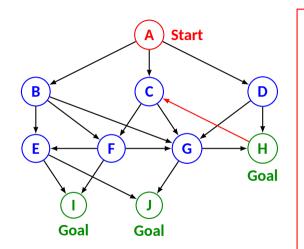
Step	<b>OPEN</b>	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB



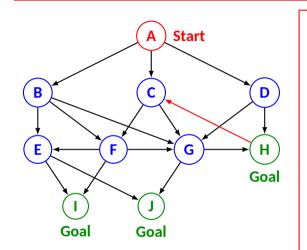
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC



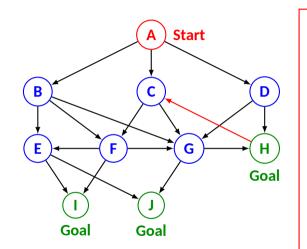
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD



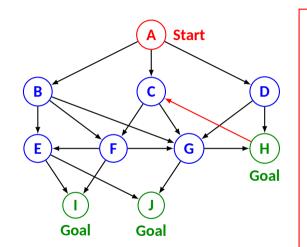
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE



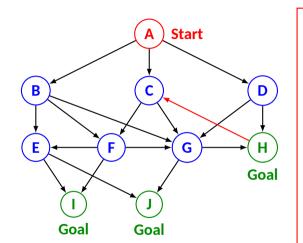
Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE
7	GHIJ	ABCDEF



Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE
7	GHIJ	ABCDEF
8	HIJ	ABCDEFG



Step	OPEN	CLOSED
1	Α	{}
2	BCD	Α
3	CDEFG	AB
4	DEFG	ABC
5	EFGH	ABCD
6	FGHIJ	ABCDE
7	GHIJ	ABCDEF
8	HIJ	ABCDEFG
9	Goal no	de H found



Step	<b>OPEN</b>	CLOSED	
1	Α	{}	
2	BCD	Α	
3	CDEFG	AB	
4	DEFG	ABC	
5	EFGH	ABCD	
6	FGHIJ	ABCDE	
7	GHIJ	ABCDEF	
8	HIJ	ABCDEFG	
9	Goal node H found		

BFS guarantees shortest length path to goal but has higher memory requirement.

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	Iterative Deepening	BFS
Complete?			

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	<b>Iterative Deepening</b>	BFS
Complete?	No	Yes	Yes

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	<b>Iterative Deepening</b>	BFS
Complete?	No	Yes	Yes
Time			

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	<b>Iterative Deepening</b>	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	<b>Iterative Deepening</b>	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space			

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	<b>Iterative Deepening</b>	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space	O(bm)	O(bd)	$O(b^d)$

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

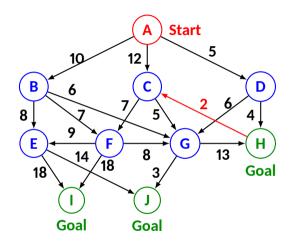
Criterion	DFS	<b>Iterative Deepening</b>	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space	O(bm)	O(bd)	$O(b^d)$
0			

**Optimal** 

- b branching factor, d depth of shallowest soln, m maximum depth
- Optimality under the assumption of identical cost for all steps

Criterion	DFS	<b>Iterative Deepening</b>	BFS
Complete?	No	Yes	Yes
Time	$O(b^m)$	$O(b^d)$	$O(b^d)$
Space	O(bm)	O(bd)	$O(b^d)$
Optimal	No	Yes	Yes

#### **Searching state space graph with edge cost**

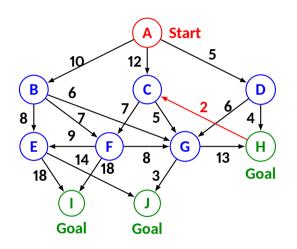


#### Modifying basic algorithms to incorporate cost

- 1. [Initialize] Initially the OPEN List contains the Start Node s. CLOSED List is Empty.
- 2. [Select] Select the first Node n on the OPEN List. If OPEN is empty, Terminate
- 3. [Goal Test] If n is Goal, then decide on Termination or Continuation / Cost Updation
- 4. [Expand]
  - a. Generate the successors  $n_1, n_2, \dots, n_k$ , of node n, based on the State Transformation Rules
  - **b.** Put n in LIST CLOSED
  - c. For each  $n_i$ , not already in OPEN or CLOSED List, put  $n_i$  in the FRONT (for DFS) / END (for BFS) of OPEN List
  - d. For each  $n_i$  already in OPEN or CLOSED decide based on cost of the paths
- 5. [Continue] Go to Step 2

Algorithm IDS Performs DFS level by level iteratively (DFS(1), DFS(2), ...and so on)

#### Searching state space graph with edge cost



#### **Cost ordered search:**

- DFBB
- Best first search
- Best first IDS
- Use of heuristic estimates: A\*, AO\*

Thank you!