## Discrete Mathematics

Graphs-II

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- The edge $u v$ completes a cycle with portion of $P$ from $v$ to $u$
- A graph $G$ is Eulerian if and only if it has at most one non-trivial component and its vertices all have even degree


## Vertex degree

- The degree of a vertex in a graph $G$, written $d_{G}(v)$ or $d(v)$, is the number of edges incident to $v$, except that each loop at $v$ counts twice.
- The maximum degree is $\Delta(G)$, the minimum is $\delta(G)$ and $G$ is regular if $\Delta(G)=\delta(G)$


## Degree-Sum formula

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- A $k$-regular graph with $n$ vertices has $\frac{n k}{2}$ edges


## Example

- For a simple graph $G$ with vertices $v_{1}, \ldots, v_{n}$ and $n \geq 3$,

$$
e(G)=\frac{\sum_{i} e\left(G-v_{i}\right)}{n-2} \text { and } d_{G}\left(v_{j}\right)=\frac{\sum_{i} e\left(G-v_{i}\right)}{n-2}-e\left(G-v_{j}\right)
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## Graphic sequences

- The degree sequence of a graph is the list of vertex degrees, usually written in non-increasing order, as $d_{1} \geq d_{2} \geq \ldots \geq d_{n}$


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- The nonnegative integers $d_{1}, d_{2}, \ldots, d_{n}$ are vertex degrees of some graph if and only if $\sum_{i} d_{i}$ is even
- For $n>1$, an integer list $d$ of size $n$ is graphic if and only if $d^{\prime}$ is graphic, where $d^{\prime}$ is obtained from $d$ by deleting its largest element $\Delta$ and subtracting 1 from its $\Delta$ next largest elements. The only 1-element graphic sequence is $d_{1}=0$.


## Exercise

- In a league with two divisions of 13 teams each, determine whether it is possible to schedule a season with each team playing nine games against teams within its division and,four games against teams in the other division.


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- Use complete graphs and counting arguments (not algebra!) to prove
- $\binom{n}{2}=\binom{k}{2}+k(n-k)+\binom{n-k}{2}$ for $0 \leq k \leq n$


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- $\binom{n}{2}=\binom{k}{2}+k(n-k)+\binom{n-k}{2}$ for $0 \leq k \leq n$
- If $\sum_{i} n_{i}=n$ then $\sum_{i}\binom{n_{i}}{2} \leq\binom{ n}{2}$


## Exercise

- Prove that the number of simple even graphs with vertex set $[n]$ is $2^{\binom{n-1}{2}}$


## Exercise

- Which of the following are graphic sequences? Provide a construction or a proof of impossibility for each.
- $5,5,4,3,2,2,2,1$
- $5,5,4,4,2,2,1,1$
- $5,5,5,3,2,2,1,1$
- $5,5,5,4,2,1,1,1$


## Exercise

- Count the cycles of length $n$ in $K_{n}$ and the cycles of length $2 n$ in $K_{n, n}$.


## Exercise

- Determine which pairs of graphs below are isomorphic.



## Exercise

- Let $G$ be the graph with vertex set $\{1, \ldots, 15\}$ in which $i$ and $j$ are adjacent if and only if their greatest common factor exceeds 1 . Count the components of $G$ and determine the maximum length of a path in $G$.


## Exercise

- Determine the values of $m$ and $n$ such that $K_{m, n}$ is Eulerian.


## Exercise

- Let $G$ be an $n$-vertex simple graph, where $n>2$. Determine the maximum possible number of edges in $G$ under each of the following conditions.
- G has an independent set of size $a$.
- G has exactly k components.
- G is disconnected.


## Exercise

- Let $G_{n}$ be the graph whose vertices are the permutations of $\{1, \ldots, n\}$, with two permutations $a_{1}, \ldots, a_{n}$ and $b_{1}, \ldots, b_{n}$ adjacent if they differ by interchanging a pair of adjacent entries ( $G_{3}$ shown below). Is $G_{n}$ connected?


Thante youl

